

Variants of Zero Forcing (ZF)

- ZF of patterns: Z, Z^-, Z_e
 - symmetry: Z_{oc}
 - ? = 0 or * : \hat{Z}
 - minor-monotone: $[Z]$
- ZF of sign patterns: Z_{\pm}
- ZF for inertias: Z_q

$$Z_+ = Z_0 \leq Z_1 \leq \dots \leq Z_n = Z$$
- ZF for Strong Arnold Property: Z_{SAP}

ZF on a graph G .

- vertices: blue or white.
 - pick some vtx blue
 - color-change rule (CCR)
 - zero forcing number
- $Z(G)$: min blue vtx that can make whole graph blue

ZF of patterns:

pattern = matrix over $\{0, *, ?\}$
nonzero \uparrow any element.

$$\begin{array}{c} \text{eq'n} \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{array}{c} x_1 \\ x_2 \\ x_3 \\ x_4 \end{array} \begin{pmatrix} 0 & * & * & 0 \\ * & * & * & 0 \\ * & * & * & * \\ 0 & 0 & * & 0 \end{pmatrix}$$

max nullity = 0

ZF: growing knowledge:

variables: blue or white
zero \uparrow uncertain.

$$\begin{array}{l} \text{eq'n 4} \Rightarrow x_3 = 0 \\ \text{eq'n 1} \Rightarrow x_2 = 0 \\ \text{eq'n 2} \Rightarrow x_1 = 0 \\ \text{eq'n 3} \Rightarrow x_4 = 0 \end{array}$$


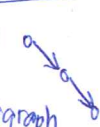


	x_1, x_2, x_3, x_4	<div style="border: 1px solid black; padding: 2px; display: inline-block;">Assume: $x_4 = 0$</div>
1	$\begin{pmatrix} 0 & * & * & 0 \\ * & * & * & 0 \\ * & * & * & * \\ 0 & 0 & * & * \end{pmatrix}$	eq'n 4 $\Rightarrow x_3 = 0$
2		eq'n 1 $\Rightarrow x_2 = 0$
3		eq'n 2 $\Rightarrow x_1 = 0$

First 3 columns are indep.

\Rightarrow max nullity $\leq \underline{1}$
 * initial blue variables.

$$\begin{array}{c} 3 \\ 4 \\ 1 \\ 2 \\ 3 \end{array} \begin{array}{c} 2 \\ 1 \\ 4 \end{array} \begin{pmatrix} * & & & * \\ * & * & & \\ * & * & * & \\ * & * & * & * \end{pmatrix}$$

Patterns / Graphs

	sym	not sym
free diagonal	$\begin{pmatrix} ? & * & 0 \\ * & ? & * \\ 0 & * & ? \end{pmatrix}$  <p>simple</p>	$\begin{pmatrix} ? & * & 0 \\ 0 & ? & * \\ 0 & 0 & ? \end{pmatrix}$  <p>simple digraph</p>
not free	$\begin{pmatrix} 0 & * & 0 \\ * & * & * \\ 0 & * & 0 \end{pmatrix}$  <p>loop</p>	$\begin{pmatrix} 0 & * & 0 \\ 0 & * & * \\ 0 & 0 & 0 \end{pmatrix}$  <p>loop digraph</p>

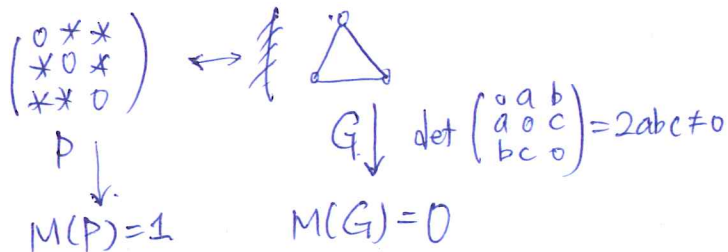
Thm: $M(G) \leq Z(G)$ [AIM'08, H'10]

BG'07
Giovannetti

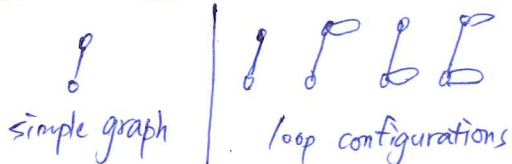
Symmetry:

G = simple or loop graph.

$M(G)$: ~~focus on symmetric mtx~~
max nullity over sym mtx of G .



? = * or 0

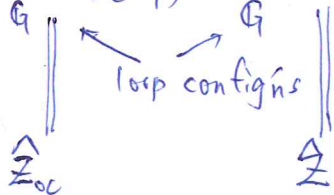


By def, $M(G) = \max_{G \leftarrow \text{all loop configs}} M(G)$

Thm: [BBFHHSvdDvdH '13, L'16]

G : simple graph

$$M(G) \leq \max_G Z_{oc}(G) \leq \max_G Z(G) \leq Z(G)$$



ZF on graphs: Z .

CCR:

simple: For vtx i , if j is the only white neighbor and i is blue

$$\Rightarrow i \rightarrow j$$

loop: For vtx i , if j is the only white neighbor

$$\Rightarrow i \rightarrow j$$

Digraph: $nbr \rightarrow out-nbr$

Odd cycle ZF: Z_{oc} (loop graph)

CCR- Z_{oc} :

• CCR- Z .

• B = blue vtx
if a component of $G \setminus B$ is a loopless odd cycle C ,
 $\Rightarrow \forall v \in V(C)$ turn blue.

Thm: [L'16]

$$M(G) \leq Z_{oc}(G) \leq Z(G)$$

G : loop graphs.

Minor-monotone:

$\beta(H) \leq \beta(G)$ if H is a minor of G .

Colin de Verdière type parameter

$$\xi(G) \leq M(G) \leq Z(G)$$

\uparrow minor-monotone

Minor-monoton floor of β

$$L\beta(G) = \min_H \beta(G) \text{ : } H \text{ is a minor of } G$$

$$\xi(G) \leq LM(G) \leq LZ(G)$$

Thm: [BBFHHSvdDvdH '13]

CCR- LZ :

• CCR- Z
• if i and all nbs are blue and i hasht make a force,
then $i \rightarrow j$ (pick a j)

ZF for sign patterns [GB'4]

Sign pattern: matrix over $\{0, ?, +, -\}$.
 any $\begin{matrix} \nearrow \\ \rightarrow \\ \searrow \end{matrix}$ $\begin{matrix} >0 <0 \\ >0 <0 \end{matrix}$

Sign ZF: Z_{\pm} .

vertices: blue $\begin{matrix} \uparrow \\ \text{zero} \end{matrix}$, white $\begin{matrix} \leftarrow \\ \uparrow \\ \text{uncertain} \end{matrix}$ $\begin{matrix} + >0 \\ - \leq 0 \end{matrix}$

CCR- Z_{\pm} :

- CCR-Z
- if all white vtxs has no mark (+, -)
pick one and ~~not~~ mark it "+"

• $\begin{matrix} \oplus \\ \oplus \\ \oplus \end{matrix} + \begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \ominus \\ \ominus \end{matrix} = 0$
 $\begin{matrix} \uparrow & \uparrow & \uparrow \\ \text{coef} & \text{var} & \text{mark it "-"} \end{matrix}$

• $\begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \ominus \\ \ominus \end{matrix} = 0$
 $\begin{matrix} \geq 0 & \geq 0 & \geq 0 \end{matrix}$

\Rightarrow all used vars turn blue

$\begin{matrix} \text{blue} & \text{blue} & \text{blue} \end{matrix}$ \leftarrow assumption.

e.g.

$\begin{matrix} + & - \\ + & - \\ + & + \end{matrix}$	$\begin{matrix} \oplus \\ \oplus \end{matrix}$	$\begin{matrix} ? \\ ? \end{matrix}$
$\begin{matrix} ? \\ ? \end{matrix}$	$\begin{matrix} \oplus \\ \ominus \\ ? \end{matrix}$	$\begin{matrix} ? \\ ? \end{matrix}$

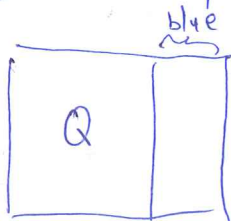
$\begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \ominus \\ \ominus \end{matrix} = 0$
 $\begin{matrix} \uparrow \\ \text{mark "+"} \end{matrix}$

$\begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \oplus \\ \oplus \end{matrix} + \begin{matrix} \oplus \\ \oplus \end{matrix} = 0$
 $\begin{matrix} \geq 0 & \geq 0 & \geq 0 \end{matrix}$

$\begin{matrix} \ominus \\ \oplus \end{matrix} + \begin{matrix} \ominus \\ \oplus \end{matrix} = 0$
 $\begin{matrix} \leq 0 & \leq 0 \end{matrix}$

Thm: max nullity \leq initial blue vars.

Question: To improve Z_{\pm} .



$\Rightarrow Q^T$ is L-matrix

L-matrix: a sign pattern s.t.h every realization has row full-rank.

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ZF for inertia: Z_q [BGH '15]

Thm: $M_q(G) \leq Z_q(G)$

G : simple graph

Consider ~~the~~ mtx with $\leq q$ neg eigenvalues.

"Game" on simple graph G .

players: X and Y

X : color all vtx blue.

Y : make X use max token

CCR- Z_q :

- X can use 1 token to color any vtx blue.

- X can apply CCR-Z on G of free

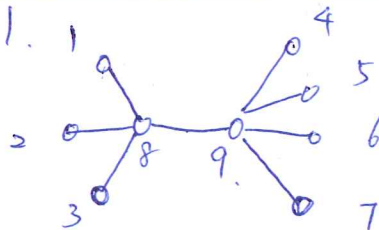
- $\begin{cases} B: \text{blue vtx} \\ G \setminus B \text{ has compnts } W_1, \dots, W_k. \end{cases}$

X can pick $\geq q+1$ compnts to Y .

Y return ≥ 1 compnts $W_{a_1}, W_{a_2}, \dots, W_{a_s}$.

X can apply CCR-Z on $G[B \cup W_{a_1} \cup W_{a_2} \dots \cup W_{a_s}]$ for free

e.g. $q=1$.



9 tokens is possible.

$N(G)$

5 tokens is possible.

$Z(G)$

4 tokens is possible.

Z_q : min tokens required

2 tokens \rightarrow blue. 3, 7

CCR-Z \rightarrow 8, 9

$X \xrightarrow{1,4} Y \xrightarrow{1} \begin{matrix} 8 & \checkmark \\ 4 & \times \\ 1,4 & \times \end{matrix}$

$X \xrightarrow{2,4} Y \xrightarrow{2} 2$

2 tokens \rightarrow 4, 5

CCR-Z \rightarrow 6

Thm: $Z_+ = Z_0 \leq Z_1 \leq Z_2 \leq \dots \leq Z_n = Z$

Question: Combine $Z_1, Z_2 \rightarrow Z_u?$

$u \leftarrow$ 1 neg eigen sgn pattern

Remark: 1. Y might not want to return all compnts

2. X might not want to spend all token initially