

# Constructing the Surreal Numbers

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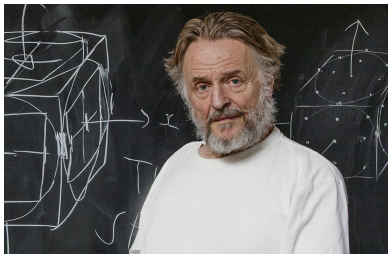
Gettysburg College

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# History of Surreal Numbers

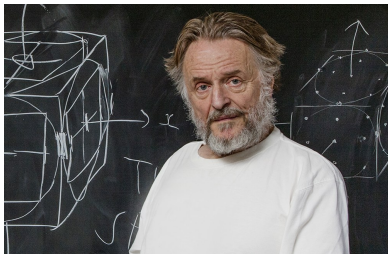
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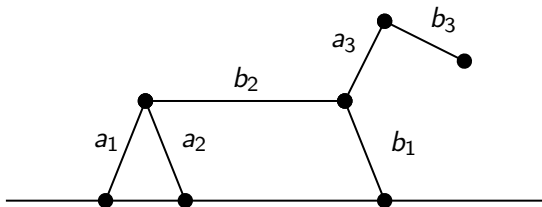
- “All Numbers Great and Small” by John Conway in 1972
- *Surreal Numbers* by Donald Knuth in 1974



# Review of *Aerion*

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Recall the following hackenbush game between players Left (who can remove pieces  $a_1, a_2, a_3$ ) and Right (who can remove pieces  $b_1, b_2, b_3$ ):



# From Hackenbush to Ordered Pairs

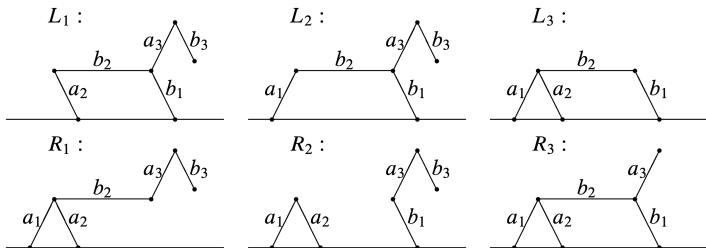
## From Hackenbush to Ordered Pairs

Let  $L_1$ ,  $L_2$ , and  $L_3$  denote the results of player Left exercising options  $a_1$ ,  $a_2$ , and  $a_3$ , respectively, and let  $R_1$ ,  $R_2$ , and  $R_3$  denote the results of player Right exercising options  $b_1$ ,  $b_2$ , and  $b_3$ , respectively.



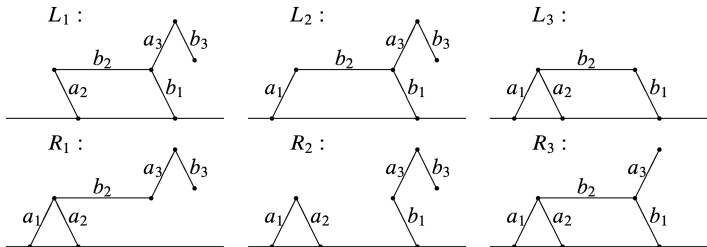
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We can represent the *Aerion* game with the ordered tuple of sets

$$\text{Aerion} = (\{L_1, L_2, L_3\}, \{R_1, R_2, R_3\}).$$

# Sequences of Games

When playing, Left and Right alternate choosing an option and thereby transform the game to another, simpler, game.

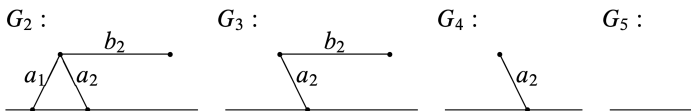
# Sequences of Games

When playing, Left and Right alternate choosing an option and thereby transform the game to another, simpler, game.

So, a particular “game-play” of Aerion is the sequence

$$G = (G_0, G_1, G_2, G_3, G_4, G_5),$$

where  $G_0 = \text{Aerion}$ ,  $G_1 = L3$ , and





## Definition

We say that a collection of ordered pairs of sets is *inductive* if it has the property that whenever  $\mathcal{L}$  and  $\mathcal{R}$  are subsets of the collection, the ordered pair  $(\mathcal{L}, \mathcal{R})$  is an element of the collection.

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## Example

Let  $\Omega$  be an inductive collection and suppose  $\{1, 2, 3\} \subset \Omega$ . Then

$$(\{1, 2\}, \{3\}) \in \Omega.$$



# Categorical Definition of $\Omega$

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Let  $\Omega$  be the (very large) collection of ordered pairs  $(\mathcal{L}, \mathcal{R})$  of sets  $\mathcal{L}$  and  $\mathcal{R}$ , that is categorically determined by the two axioms:

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## Axioms

- **Axiom 24.1** The collection  $\Omega$  is inductive.
- **Axiom 24.2** The only inductive subcollection of  $\Omega$  is  $\Omega$  itself.

# Building Games

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Thus we have that  $\Omega_1 = \{G(1), G(-1), N(1)\}$ .

We can continue, but it turns out that  $|\Omega_2| = 252$ , and clearly, generations continue to grow in size.

# An Order Relation on $\Omega$

## Definition 24.3.

Let  $G = (\mathcal{L}_G, \mathcal{R}_G) \in \Omega$  and  $H = (\mathcal{L}_H, \mathcal{R}_H) \in \Omega$ . We say that  $G$  is “less than or similar” to  $H$  and write  $G \lesssim H$  if, and only if:

- There is no  $L \in \mathcal{L}_G$  for which  $H \lesssim L$ .
- There is no  $R \in \mathcal{R}_H$  for which  $R \lesssim G$ .

If  $G$  is not less than or similar to  $H$ , we write  $G \not\lesssim H$ .

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## Proposition 24.4.

For any  $\mathcal{X}, \mathcal{Y} \in \Omega$ , we have that  $(\emptyset, \mathcal{X}) \lesssim (\mathcal{Y}, \emptyset)$ .

# Addition in $\Omega$

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Now we define an addition operation on  $\Omega$ .



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## Definition 24.22.

Let  $G = (\mathcal{L}_G, \mathcal{R}_G) \in \Omega$  and  $H = (\mathcal{L}_H, \mathcal{R}_H) \in \Omega$ . We define the sum of  $G$  and  $H$  recursively as

$$G + H = ((G + \mathcal{L}_H) \cup (\mathcal{L}_G + H), (G + \mathcal{R}_H) \cup (\mathcal{R}_G + H)).$$

Here, as customary, the sum of a game and a collection of games is simply the collection of the appropriate sums; for example,

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## Theorem 24.23.

Addition on  $\Omega$  is a closed operation, is commutative, is associative, has the identity property, and has the additive inverse property (up to similarity).

# Multiplication in $\Omega$

We also define a multiplication operation on  $\Omega$ .

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## Definition 24.26.

Let  $G = (\mathcal{L}_G, \mathcal{R}_G) \in \Omega$  and  $H = (\mathcal{L}_H, \mathcal{R}_H) \in \Omega$ . We define the product of  $G$  and  $H$  recursively as

$$G \cdot H = ((G \cdot \mathcal{L}_H + \mathcal{L}_G \cdot H - \mathcal{L}_G \cdot \mathcal{L}_H) \cup (G \cdot \mathcal{R}_H + \mathcal{R}_G \cdot H - \mathcal{R}_G \cdot \mathcal{R}_H), \\ (G \cdot \mathcal{R}_H + \mathcal{L}_G \cdot H - \mathcal{L}_G \cdot \mathcal{R}_H) \cup (G \cdot \mathcal{L}_H + \mathcal{R}_G \cdot H - \mathcal{R}_G \cdot \mathcal{L}_H)).$$



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## Definition 24.7.

The collection of numeric games  $\Gamma$  consists of games  $(\mathcal{L}, \mathcal{R})$  where  $\mathcal{L} \subseteq \Gamma$ ,  $\mathcal{R} \subseteq \Gamma$ , and there is no  $L \in \mathcal{L}$  and  $R \in \mathcal{R}$  for which  $R \lesssim L$ .



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## Proposition 24.8.

For any  $\mathcal{X} \in \Gamma$ , we have  $(\emptyset, \mathcal{X}) \in \Gamma$  and  $(\mathcal{X}, \emptyset) \in \Gamma$ .

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Write  $\Gamma_n = \Gamma \cap \Omega_n$ .

# More Numeric Games

Using Definition 24.7 we can get

$$\Gamma_0 = \{G(0)\}, \text{ and } \Gamma_1 = \{G(-1), G(1)\}.$$

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It turns out that  $|\Gamma_2| = 17$ . Here are four elements of  $\Gamma_2$ :

$$G(-2) = (\emptyset, \{G(-1), G(0)\});$$

$$G\left(-\frac{1}{2}\right) = (\{G(-1)\}, \{G(0)\});$$

$$G\left(\frac{1}{2}\right) = (\{G(0)\}, \{G(1)\});$$

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We will see shortly why we only have to consider these four!

# Extending the Relation on $\Omega$

## Definition 24.5.

Suppose that  $G \in \Omega$  and  $H \in \Omega$ . We say that:

- $G$  is “similar” to  $H$  and write  $G \approx H$ , if  $G \lesssim H$  and  $H \lesssim G$ .
- $G$  is “greater than” or similar to  $H$  and write  $G \gtrsim H$  if  $H \lesssim G$ .
- $G$  is “less than”  $H$  and write  $G < H$ , if  $G \lesssim H$  and  $G \not\approx H$ .
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## Proposition 24.12 and 24.14.

Let  $G$ ,  $H$ , and  $K$  be arbitrary games, and suppose that  $G \lesssim H$  and  $H \lesssim K$ . Then we also have  $G \lesssim G$  and  $G \lesssim K$ .

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## Corollary 24.20.

The relation  $\approx$  of similarity is an equivalence relation on the collection of games  $\Omega$ .

# Defining Surreal Numbers

The similarity relation partitions  $\Omega$  into equivalence classes, but we are only interested in the equivalence classes within  $\Gamma$ .

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We denote the equivalence class of a game  $G$  by  $[G]$ ; that is,

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## Definition 24.21.

The equivalence classes of the similarity relation in  $\Gamma$  are called *surreal numbers*. The collection of surreal numbers is denoted by  $\mathbb{S}$ .

# An Example of Equivalence Relation (Proposition 24.9)

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$G$	$G \in \dots$	$\dots \in [G]$
$G(0)$	$\Gamma_0$	$(\{G(-1)\}, \emptyset),$ $(\emptyset, \{G(1)\}),$ $(\{G(-1)\}, \{G(1)\})$
$G(1)$	$\Gamma_1$	$(\{G(-1), G(0)\}, \emptyset)$
$G(-1)$	$\Gamma_1$	$(\emptyset, \{G(0), G(1)\})$
$G(2)$	$\Gamma_2$	$(\{G(1)\}, \emptyset, \emptyset)$ $(\{G(-1), G(1)\}, \emptyset),$ $(\{G(-1), G(0), G(1)\}, \emptyset)$
$G(-2)$	$\Gamma_2$	$(\emptyset, \{G(-1)\})$ $(\emptyset, \{G(-1), G(1)\}),$ $(\emptyset, \{G(-1), G(0), G(1)\})$
$G(1/2)$	$\Gamma_2$	$(\{G(-1), G(0)\}, \{G(1)\})$
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# $\mathbb{S}$ is an ordered field

## Definition 24.27.

Let  $x_{\mathbb{S}}$  and  $y_{\mathbb{S}}$  be surreal numbers, and suppose that  $G \in x_{\mathbb{S}}$  and  $H \in y_{\mathbb{S}}$ . We define surreal addition and multiplication as

- $x_{\mathbb{S}} + y_{\mathbb{S}} = [G + H]$ ; and
- $x_{\mathbb{S}} \cdot y_{\mathbb{S}} = [G \cdot H]$ .

Furthermore, we say that  $x_{\mathbb{S}}$  is positive if  $G > G(0)$  and that it is negative if  $G < G(0)$ ; we let  $P_{\mathbb{S}}$  denote the collection of positive surreal numbers.

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Here is the main result:

## Theorem 24.28.

The surreal number system  $(\mathbb{S}, +, \cdot, P_{\mathbb{S}})$  is an ordered field.

# Interesting Results

## Theorem 24.29.

The surreal number system  $(\mathbb{S}, +, \cdot, P_{\mathbb{S}})$  is the largest ordered field; that is, for every ordered field system  $(F, +_F, \cdot_F, P_F)$ , there exist a subfield  $S \subseteq \mathbb{S}$  for which the systems  $(F, +_F, \cdot_F, P_F)$  and  $(S, +_S, \cdot, P_S)$  are isomorphic.

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As a Corollary to Axiom 24.2 we have,

## Theorem 24.10. (Conway's Induction Principle)

Suppose that  $P(G)$  is a predicate that becomes a statement for all  $G \in \Omega$ . If the implication

$$\left( \bigwedge_{K \in \mathcal{L} \cup \mathcal{R}} P(K) \right) \implies P(H)$$

holds for all  $H = (\mathcal{L}, \mathcal{R}) \in \Omega$ , then  $P(G)$  is true for every  $G \in \Omega$ .